#### Applications

# Indefinites: Exceptional Scope and Marked Indefinites

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Team Semantics and Dependence 13 January 2022

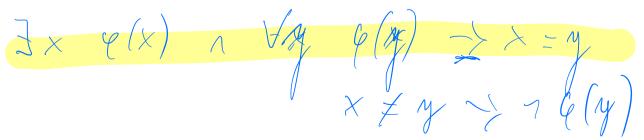
# Plan for today

- Remarks on yesterday's exercise
- What is an indefinite?
- Indefinites and exceptional scope
- Marked indefinites
  - Desiderata
  - Two-sorted team semantics
  - Applications

Yesterday's exercise

- a. Exactly one student can take Spanish or Calculus. (1)
  - b. ---- One student can choose between the two and each of the others can take neither of them.
- a. Exactly one student cannot take Spanish or Calculus. (2)
  - b. ---- One student can take neither Calculus nor Spanish and each of the others can choose between them.

Different definitions of implication are possible in state-based systems. In inquisitive semantics:  $M, s \models A \rightarrow B$  iff for all  $t \subseteq s : M, t \models A \Rightarrow M, t \models B$ 



# What is an indefinite? (1)

- (1) a. Sue likes a book, some book, a certain book.
  - b. Sue likes this book, these books, the book.

(1a) contains examples of determiner phrases which are headed by an **indefinite** determiner.

(1b) contains examples of determiner phrases which are headed by an **definite** determiner.

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(1a) contains examples of determiner phrases which are headed by an **indefinite** determiner.

(1b) contains examples of determiner phrases which are headed by an **definite** determiner.

(1a) and (1b) are treated as existentials as opposed to universal DPs (e.g. *every book*).

# What is an indefinite? (2)

Determining what counts as an indefinite has led to important turning points in formal semantics.

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In the classical *generalized quantifier theory* (Montague 1973, Barwise and Cooper 1981, Keenan and Stavi 1986), indefinites are a subtype of GQ which are **non-unique**, as opposed to definites, which are **unique**:

(2) a. a book 
$$\mapsto \lambda P \exists (B(x) \land P(x))$$

b. the book  $\mapsto \lambda P \exists x \forall \mathbf{y}((\mathbf{B}(\mathbf{y}) \leftrightarrow \mathbf{x} = \mathbf{y}) \land P(x))$ 

# What is an indefinite? (3)

In the dynamic tradition (Karttunen 1976, Kamp 1981, Heim 1982, Dekker 1993), indefinites introduce new discourse variables. Indefinites are thus **novel**, while definites are **familiar**:

- (3) a. A book<sub>x</sub> is on the table. It<sub>x</sub>/The book<sub>x/y</sub>/A book<sub>#x/y</sub> is called 'War and Peace'.
  - b. Every  $book_{x/1}$  on the table.  $It_{\#x/y}$  is called 'War and Peace'.

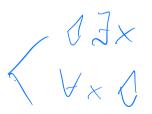
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  - b. Every book<sub>x</sub> on the table.  $It_{\#x/y}$  is called 'War and Peace'.

What about any book or no book?

- (4) a. You can take any book.
  - b. You read no book.



No unique answer: *any book* can be an indefinite existential interpreted within the scope of the modal *can*; or a universal with a modal in its scope. Similarly, *no book* can be an existential within the scope of a covert negation; or a negative universal quantifier.

# Indefinites and Freedom of Scope

A salient property of indefinites is their ability to take **scope freely** over several operators:

- (5) a. Sue likes every book which concerns an important war.  $\forall \times \exists \psi ; \exists \psi \forall \times$ 
  - b. Sue likes a book which concerns every important war.
- (6) a. If a panda comes to the party, Kola the bear will be happy.
  - b. If every panda comes to the party, Kola the bear will be happy.

In (5a) and (6a), the indefinite can take scope freely (even outside its syntactic boundaries). By contrast, universals are clause bound.

#### Exceptional Scope: Ambiguity Thesis Fodor & Sag (1982) treat wide-scope indefinites as referring expressions (e.g. like a proper name):

- (7) If a panda comes to the party, Kola the bear will be happy.
  - a.  $\exists x(P(x) \land C(x, p)) \rightarrow H(k)$  [quantificational]
  - b.  $\mathbf{P}(\mathbf{x}) \land C(x, p) \rightarrow H(k)$  [referential]

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Indefinites are **ambiguous** between a referential and a quantificational reading, as opposed to universals.

But what about intermediate readings (Farkas 1981)?

- (8) Every<sub>x</sub> student read every<sub>y</sub> paper that  $a_z$  professor recommended.
  - a. Narrow Scope (NS):  $\forall x/\forall y/\exists z$
  - b. Intermediate Scope (IS):  $\forall x / \exists z / \forall y$
  - c. Wide Scope (WS):  $\exists z/\forall x/\forall y$

### **Choice Functional Approaches**

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- (9) Every<sub>x</sub> student read every<sub>y</sub> paper that  $a_z$  professor recommended.
  - a. Narrow Scope (NS):  $\forall x \forall y \exists f((S(x) \land A(y) \land W(f(P), y)) \rightarrow R(x, y))$
  - b. Intermediate Scope (IS):  $\forall x \exists f \forall y ((S(x) \land A(y) \land W(f(P), y)) \rightarrow R(x, y))$
  - c. Wide Scope (WS):  $\exists f \forall x \forall y ((S(x) \land A(y) \land W(f(P), y)) \rightarrow R(x, y))$

# Other accounts

- Kratzer (1998), Matthewson (1999): the value of the choice function is unique (as in Fodor & Sag); intermediate readings are obtained by relativizing the choice function to other variables (e.g. a Skolem function  $f_x$ ).
- Abusch (1994): indefinites analyzed enter the semantic composition in a free way (implemented via a quantifier Cooper storage mechanism)
- Schwarzschild (2002): Indefinites are quantificational existentials. Exceptional scope is obtained by pragmatic restriction of the denotation of the existential to a singleton.
- Charlow (2019): Alternative Semantics analysis of indefinites and scope taking.

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The interpretation function is of the form  $[[-]]^{M,G,V}$ , where G is a **set of assignments** and V the set of variables introduced by previous operators.

- (10) Every<sub>x</sub> student read every<sub>y</sub> paper that  $a_z$  professor recommended.
  - a. Narrow Scope (NS):  $\forall x \forall y \exists^{x,y} z \phi(x, y, z)$
  - b. Intermediate Scope (IS):  $\forall x \forall y \exists^{x} z \phi(x, y, z)$
  - c. Wide Scope (WS):  $\forall x \forall y \exists^{\emptyset} z \phi(x, y, z)$

 $\exists^{U}z$  means the values of z are (possibly) different for any different value assigned to the variables in U. If  $U = \emptyset$ , then the choice of values for z is independent.



# How do Brasoveanu and Farkas (2011) deal with binder roof configurations like (11) below?

(11) Every  $man_x$  read a paper<sub>y</sub> which  $he_x$  had written.

# Varieties of Indefinites

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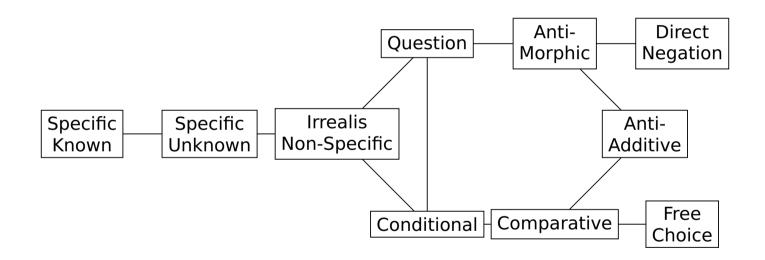
Indefinites constitute a **functionally rich** linguistic environment and display a significant cross-linguistic variety (Haspelmath 1997):

- English: *some*, *any*, *no*, ...
- Italian: qualcuno, qualunque, nessuno, (un) qualche, ...
- Dutch: iets, enig, wie dan ook, niets, ...
- Russian: koe-, -to, -nibud, ni-, ...

• . . .



# Haspelmath (1997) proposed a map capturing the functional distribution of indefinites:



We distinguish between indefinite pronouns (or determiners) like *someone* and lexical indefinite expressions like *a book*.

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#### Preliminaries

Here we focus on the *specific known*, *specific unknown* and *non-specific* functions:

- (12) a. Specific known: Someone called. I know who.
  - b. Specific unknown: Someone called. I do not know who.
  - c. Non-specific: John wants to talk with someone.

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Possible marked indefinites based on these functions:

	functions			ovampla
type	sk	su	ns	example
(i) unmarked	$\checkmark$	$\checkmark$	$\checkmark$	Italian <i>qualcuno</i>
(ii) specific	$\checkmark$	$\checkmark$	X	Georgian <i>-ghats</i>
(iii) non-specific	X	×	1	Russian <i>-nibud</i>
(iv) epistemic	//	$\checkmark$	the V	German <i><sup></sup>irgend</i> -
(v) specific known	$\checkmark$	×	X	Russian <i>‡koe</i> 👡
(vi) SK + NS	$\checkmark$	×	$\checkmark$	unattested
(vii) specific unknown	X	$\checkmark$	X	Kannada -oo 1/4

### **Our Goals**

We develop a two-sorted team semantics which accounts for:

- (a) the specific known, specific unknown and non-specific readings;
- (b) the variety of marked indefinites mentioned before;
- (c) the licensing of non-specific indefinites;
- (d) the relationship between scope and marked indefinites;
- (e) the diachronic pathway from non-specific to epistemic;
- (f) the fine-grained analysis of epistemic indefinites (*irgend-*).

# Break



Keukenhof, Lisse

# Licensing of non-specific indefinites

Non-specific indefinites are **ungrammatical in episodic sentences** and they need an operator (e.g. a universal quantifier or a modal) which licenses them:

(13) \**Ivan včera kupil kakuju-nibud' knigu.* Ivan yesterday bought which-indef. book.

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Haspelmath (1997) proposed that indefinites might change due to weakening of functions from the right (non-specific) of the functional map to the left (specific).

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(15) Weakening of functions (a) specific known < (b) specific unknown < (c) non-specific

But then why diachronically we do not observe the acquisition of (a)?

#### Interaction with scope

Specific indefinites have only wide-scope readings, while indefinites which admit non-specific readings allow for all scope configurations:

Yx Yy Jx

	WS	IS	NS
unmarked	$\checkmark$	$\checkmark$	$\checkmark$
specific	$\checkmark$	X	X
non-specific	×	$\checkmark$	$\checkmark$
epistemic	$\checkmark$	$\checkmark$	$\checkmark$
specific known	$\checkmark$	X	X
specific unknown	$\checkmark$	X	X

### **Epistemic Indefinites (Irgendein)**

*Irgend*- indefinites, unlike other epistemic indefinites (e.g. Spanish algún or Italian un gualche) admit also free choice readings:

- Mary musste irgendeinen Mann heiraten. (16)Mary had-to irgend-one man marry.
  - a. Specific unknown: There was some man Mary had to marry, the speaker doesn't know or care who it was.
  - b. Free choice: Mary had to marry a man, any man was a permitted marriage option for her.

	episodic	epistemic modal	root modal
specific unknown	1	$\checkmark$	$\checkmark$
non-specific	X	$\checkmark$	$\checkmark$
free choice	X	<b>√</b> -X	$\checkmark$

#### The Framework: Language & Team

#### **Language:** t ::= c | v $\phi ::= P(\vec{t}) | \phi \lor \psi | \phi \land \psi | \exists v \phi | \forall v \phi$

#### Team:

Given a first-order model  $M = \langle D, I \rangle$  and a sequence of variables  $\vec{v}$ , a team T over M with domain  $Dom(T) = \vec{v}$  is a set of variable assignments from  $\vec{v}$  to Dom(M) = D.

#### The Framework: Semantic Clauses

$$M,T \models P(x_1,\ldots,x_n)$$

$$M,T\models\phi\wedge\psi$$

$$M,T\models\phi\vee\psi$$

 $M, T \models \forall v \phi$ 

 $M,T \models \exists_{\text{strict}} y \phi$ 

 $M, T \models \exists_{\text{lax}} y \phi$ 

 $\Leftrightarrow \quad \forall i \in T : \langle [x_1]_{M,i}, \dots, [x_n]_{M,i} \rangle \in$  $I(P^n)$ 

$$\Leftrightarrow M, T \models \phi \text{ and } M, T \models \psi$$

there is a team  $T = T_1 \cup T_2$  s.t.  $\Leftrightarrow$  $M, T_1 \models \phi$  and  $M, T_2 \models \psi$ 

$$\Leftrightarrow M, T[/y] \models \phi, \text{ where } T[/y] = \{i[x/d] | i \in T \text{ and } d \in D\}$$

- there is a function  $h : T \rightarrow D$  $\Leftrightarrow$ s.t.  $M, T[h/y] \models \phi$ , where  $T[h/y] = \{i[h(i)/y] : i \in T\}$
- $\Leftrightarrow$  there is a function f :  $T \rightarrow$  $\wp(D) \setminus \{\emptyset\}$  s.t.  $M, T[f/y] \models \phi$ , where  $T[f/y] = \{i[d/y] : i \in$  $T, d \in f(i)$

 $S = 2 d_1 d_1$ 

Illustrations (1)

## $M, T \models \forall y \phi \Leftrightarrow M, T[/y] \models \phi$ , where $T[/y] = \{i[x/d] | i \in T \text{ and } d \in D\}$

T	X	<i>T</i> [/y]	x	У
<i>i</i> 1	$d_1$	<i>i</i> <sub>11</sub>	$d_1$	$d_1$
i2	d <sub>2</sub>	i <sub>12</sub>	$d_1$	d <sub>2</sub>
		i <sub>21</sub>	$d_2$	$d_1$
		i <sub>22</sub>	$d_2$	$d_2$

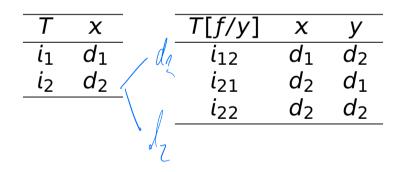
#### Illustrations (2)

## $\begin{array}{l} M,T \models \exists_{\text{strict}} y \phi \Leftrightarrow \text{ there is a function } h: T \to D \text{ s.t.} \\ M,T[h/y] \models \phi, \text{ where } T[h/y] = \{i[h(i)/y]: i \in T\} \end{array}$

Т	X	T[h/y]	X	У
<i>i</i> 1	$d_1$	i <sub>12</sub>	$d_1$	d2
i2	<i>d</i> <sub>2</sub>	i <sub>21</sub>	<i>d</i> <sub>2</sub>	$d_1$

Illustrations (3)

# $\begin{array}{l} M,T \models \exists_{lax} y \phi \Leftrightarrow \text{ there is a function } f:T \rightarrow \wp(D) \setminus \{\emptyset\} \text{ s.t.} \\ M,T[f/y] \models \phi, \text{ where } T[f/y] = \{i[d/y]: i \in T, d \in f(i)\} \end{array}$



#### Dependence Atoms

Dependence atoms impose conditions of dependence on the variable's values across different assignments:

#### **Dependence Atom:**

$$M, T \models dep(\vec{x}, \vec{y}) \Leftrightarrow \text{ for all } i, j \in T : i(\vec{x}) = j(\vec{x}) \Rightarrow i(\vec{y}) = j(\vec{y})$$

#### **Variation Atom:**

 $M, T \models var(\vec{x}, \vec{y}) \Leftrightarrow \text{ there is } i, j \in T : i(\vec{x}) = j(\vec{x}) \& i(\vec{y}) \neq j(\vec{y})$ 

Т	X	У	Ζ	l
i	$a_1$	$b_1$	<i>C</i> 1	$d_1$
j	$a_1$	$b_1$	<i>C</i> <sub>2</sub>	$d_1$
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Т	X	У	Ζ	l
i	$a_1$	$b_1$	<i>C</i> 1	$d_1$
j	$a_1$	$b_1$	<i>C</i> <sub>2</sub>	$d_1$
k	a <sub>3</sub>	b <sub>2</sub>	С3	$d_1$

Is it always the case that  $\phi \lor \phi$  is equivalent to  $\phi$ ?

 $dep(\mathcal{P}_{i} \times) \vee dep(\mathcal{P}_{i} \times)$  $\exists \times \_ depters$ 

#### **Exceptional Scope**

Similarly to B&F (2011), we interpret indefinites *in-situ* and we assume that an indefinite  $\exists x$  in syntactic scope of  $O_{\vec{z}}$  allows all  $dep(\vec{y}, x)$ , with  $\vec{y}$  included in  $\vec{z}$ .

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- (17) Every student<sub>x</sub> read every  $book_z$  that a professor<sub>y</sub> recommended.
  - a. WS  $[\exists y/\forall x/\forall z]$ :  $\forall x \forall z \exists y (\phi \land dep(\emptyset, y))$
  - b. NS  $[\forall x/\forall z/\exists y]$ :  $\forall x\forall z\exists y(\phi \land dep(xz, y))$
  - c. IS  $[\forall x/\exists y/\forall z]$ :  $\forall x\forall z\exists y(\phi \land dep(x, y))$

X	Z	y y		X	Z	y y		X	Z	<b>y</b>
		<i>b</i> <sub>1</sub>	-	<i>a</i> <sub>1</sub>	<i>C</i> <sub>1</sub>	$b_1$	_	$a_1$		$b_1$
		b <sub>1</sub>		$a_1$	<i>C</i> <sub>2</sub>	b2		$a_1$		$b_1$
					<i>C</i> <sub>1</sub>					
		$b_1$	-	a2	<i>C</i> <sub>2</sub>	$b_4$		a2	• • •	b <sub>2</sub>

WS:  $dep(\emptyset, y)$ 

NS: dep(xz, y)

IS: dep(x, y)

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We use a **two-sorted** framework with v as variable for the actual world: in the specific known, the referent is constant across all epistemically possible worlds; in the specific unknown it will vary across epistemically possible worlds.

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**Initial Team:** A team T is *initial* iff  $Dom(T) = \{v\}$ .

A sentence is **felicitous/grammatical** if there is an initial team which supports it.

#### Specific Known, Specific Unknown, Non-specific

$v$ $v$ $x$ constancy $dep(\emptyset, x)$ $\dots$ $d_1$ $v$ $v$ $d_1$ $\dots$ $v$ $v$ $x$ $\dots$ $v$ $u$ $d_1$ $\dots$ $v$ $v$ $u$ $d_2$ $v$ -constancy $dep(v, x)$ $w_1$ $d_1$ $w_2$ $d_2$ $v$ $x$ $v$ -variation $var(v, x)$ $w_1$ $d_1$ $w_1$ $d_2$ $w_1$ $d_1$				
$\dots$ $d_1$ variation $var(\emptyset, x)$ $\begin{matrix} v & x \\ \dots & d_1 \\ \hline & & \ddots & d_1 \\ \hline & & & & \ddots \\ \hline & & & & & & \\ \hline v & -vonstancy$ $war(\emptyset, x)$ $v$ $v$ $x$ $v$ -constancy $dep(v, x)$ $\begin{matrix} v & x \\ \hline w_1 & d_1 \\ \hline w_2 & d_2 \\ \hline v & x \\ \hline w_1 & d_1 \\ \hline w_1 & w_1 \\ \hline w_1 & d_1 \\ \hline w_1 & w_1 \\ \hline w_$			ν	x
variationvar(Ø, x) $v$ v-constancy $dep(v, x)$ $w_1$ v-constancy $dep(v, x)$ $w_1$ $w_1$ $d_1$ $w_2$ $d_2$ v-variation $var(v, x)$ $w_1$	constancy	dep(Ø, x)		$d_1$
variation $var(\emptyset, x)$ $\dots$ $d_1$ $\dots$ $d_2$ $\nu$ -constancy $dep(\nu, x)$ $\frac{\nu}{w_1}$ $d_1$ $w_2$ $d_2$ $\nu$ $x$ $\nu$ -variation $var(\nu, x)$ $w_1$ $d_1$				<i>d</i> <sub>1</sub>
$v-constancy  dep(v,x)  \frac{v  x}{w_1  d_1} \\ w_2  d_2 \\ \hline v-variation  var(v,x)  \frac{w_1  d_1}{w_1  d_1} \\ \hline w_1  d_1 \\ \hline w_1 \\ $	variation		ν	X
$\begin{array}{c c} v & constancy & dep(v,x) & \hline v & x \\ \hline w_1 & d_1 \\ \hline w_2 & d_2 \\ \hline v & v \\ \hline v \\ v \\ \hline v \\ v \\ \hline w_1 & d_1 \\ \hline w_1 & d_1 \\ \hline \end{array}$		var(Ø, x)		$d_1$
$v$ -constancy $dep(v, x)$ $w_1$ $d_1$ $w_2$ $d_2$ $v$ -variation $var(v, x)$ $w_1$ $d_1$				d <sub>2</sub>
$w_2  d_2$ $v_2  v_2$ $v_1  d_1$			ν	Х
v-variation $var(v, x) = \frac{v + x}{w_1 + d_1}$	v-constancy	dep(v, x)	$w_1$	$d_1$
v-variation $var(v, x)$ $w_1 d_1$			W <sub>2</sub>	d <sub>2</sub>
	u-variation		ν	x
$w_1  d_2$		var(v,x)	$w_1$	$d_1$
			$w_1$	$d_2$

Specific Known: constancy

**Specific Unknown:** *v*-constancy + variation

**Non-specific:** v-variation

Marked Indefinites

nites Desiderata

### Variety of Indefinites

type	functions			requirement	example	
type	sk	su	ns	requirement	example	
(i) unmarked	$\checkmark$	$\checkmark$	$\checkmark$	none	Italian <i>qualcuno</i>	
(ii) specific	$\checkmark$	$\checkmark$	X	dep(v, x)	Georgian <i>-ghats</i>	
(iii) non-specific	X	X	$\checkmark$	var(v,x)	Russian <i>=nibud</i>	
(iv) epistemic	$\checkmark$	$\checkmark$	X	var(Ø, x)	German <i>yjrgend</i>	
(v) specific known	$\checkmark$	X	X	dep(Ø,x)	Russian -koe	
(vi) SK + NS	$\checkmark$	X	$\checkmark$	$dep(\emptyset, x) \lor var(v, x)$	unattested	
(vii) specific unknown	X	1	X	$dep(v, x) \land var(\emptyset, x)$	Kannada <i>-oo</i>	

(vi) SK + NS: violation of connectedness (Gardenfors 2014; Enguehard and Chemla 2021)

#### (vii) specific unknown: increased complexity

(We will revise these requirements in light of our discussion about scope)

Non-specific indefinites cannot occur freely in episodic sentences, but they need an operator to be licensed.

Recall that non-specific indefinites trigger v-variation: var(v, x).

ν  $W_1$  $W_2$ 

Non-specific indefinites cannot occur freely in episodic sentences, but they need an operator to be licensed.

Recall that non-specific indefinites trigger v-variation: var(v, x).

 $\exists x \phi$ 

ν	ν	X
$W_1$	$w_1$	$a_1$
W <sub>2</sub>	W <sub>2</sub>	a2

Non-specific indefinites cannot occur freely in episodic sentences, but they need an operator to be licensed.

Recall that non-specific indefinites trigger v-variation: var(v, x).

 $\forall y \phi$ 

ν	ν	У
$W_1$	$w_1$	$b_1$
W <sub>2</sub>	$w_1$	b <sub>2</sub>
	W <sub>2</sub>	$b_1$
	W <sub>2</sub>	b <sub>2</sub>

Non-specific indefinites cannot occur freely in episodic sentences, but they need an operator to be licensed.

Recall that non-specific indefinites trigger v-variation: var(v, x).

 $\forall y \exists x \phi$ 

$\overline{\nu}$	ν	У	-	ν	У	X
$w_1$	$w_1$	$b_1$	-	$w_1$	$b_1$	$a_1$
W <sub>2</sub>	$w_1$	b <sub>2</sub>		$w_1$	b <sub>2</sub>	a2
	W <sub>2</sub>	$b_1$		W <sub>2</sub>	$b_1$	$a_1$
	<i>W</i> <sub>2</sub>	b <sub>2</sub>		W <sub>2</sub>	b <sub>2</sub>	a2

Non-specific indefinites cannot occur freely in episodic sentences, but they need an operator to be licensed.

Recall that non-specific indefinites trigger v-variation: var(v, x).

#### $\forall y \exists x \phi$

ν	ν	У	-	ν	У	x
$w_1$	$w_1$	$b_1$	-	$w_1$	$b_1$	<i>a</i> <sub>1</sub>
W2	$w_1$	b <sub>2</sub>		$w_1$	b <sub>2</sub>	a <sub>2</sub>
	W <sub>2</sub>	$b_1$		W <sub>2</sub>	$b_1$	$a_1$
	W <sub>2</sub>	b <sub>2</sub>		W <sub>2</sub>	b <sub>2</sub>	a <sub>2</sub>

Note: Indefinites can also be licensed by modals, which are analyzed via the lax existential.

#### Applications

#### Interaction with Scope

We have assumed that marked indefinites trigger the activation of particular atoms. But dependency atoms also account for scope.

**Plain:**  $dep(\vec{y}, x)$ Specific:  $dep(\vec{y}, x)$  with  $\vec{y} \subseteq \{v\}$   $\forall \vec{z} = \emptyset$ ,  $h_{y}(\vec{y}, x)$ **Epistemic:**  $dep(\vec{y}, x) \wedge var(\vec{z}, x)$  with  $\vec{z} \subseteq Var(W)$ **Non-specific:**  $dep(\vec{y}, x) \wedge var(\vec{z}, x)$  with  $\vec{z} \subseteq Var(W)$  and ź≠Ø

#### Illustration

#### $\forall z \forall y \exists z \phi$

	wide scope <i>dep</i> (v,x)	intermediate scope dep(vyz, x)	narrow scope dep(vy, x)
unmarked	$\checkmark$	$\checkmark$	$\checkmark$
specific $dep(v, x)$	1	×	×
non-specific var(v,x)	×	$\checkmark$	$\checkmark$
epistemic var(Ø, x)	$\checkmark$	1	$\checkmark$
specific known <i>dep</i> (Ø, x)	$\checkmark$	×	×
specific unknown $dep(v, x) \land var(\emptyset, x)$	1	×	×

#### (18) Weakening of functions

(a) specific known < (b) specific unknown < (c) non-specific</li>

This framework makes the notion of weakening precise in terms of **logical entailment**.

We have weakening from non-specific to epistemic:  $var(\emptyset, x)$  implies var(y, x), but no further weakening triggering the acquisition of SK.

### Epistemic Indefinites (1)

We generalize the variation atom to express the cardinality of the variation and splitting:

 $M,T \models var_n(\vec{y},x)$  iff  $\forall d \in D^* \subseteq D$  with  $|D^*| \ge n$ , for all  $i \in I$ T, there is a  $j \in T_{i,\vec{y}}$  s.t. j(x) = d, where  $T_{i,\vec{y}} = \{j \in T | i(\vec{y}) = i\}$  $i(\vec{y})$ 

### Epistemic Indefinites (1)

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We assume that *irgend* associate with  $var_2(\emptyset, x)$ . Non-specific readings are obtained via dependency atoms assuming a domain  $\ge 2$ .

 $var_{|D|}(x, v)$  models **free choice** (full non-specificity), triggered by prosodic prominence.

#### **Epistemic Indefinites (2)**

- (19) Jeder<sub>v</sub> Student hat irgendein<sub>x</sub> Buch gelesen. every student has irgendein book read.
  - a. specific unknown:  $\forall y \exists x (\phi \land dep(v, x) \land var_2(\emptyset, x))$
  - b. non-specific:  $\overline{\forall y} \exists x (\phi \land dep(vy, x) \land var_2(\emptyset, x))$

## Epistemic Indefinites (2)

- (19) Jeder<sub>y</sub> Student hat  $irgendein_x$  Buch gelesen. every student has irgendein book read.
  - a. specific unknown:  $\forall y \exists x (\phi \land dep(v, x) \land var_2(\emptyset, x))$
  - b. non-specific:  $\forall y \exists x (\phi \land dep(vy, x) \land var_2(\emptyset, x))$
- (20) *Mary musste<sub>y</sub> irgendeinen Mann heiraten.* Mary had-to irgend-one man marry.
  - a. specific unknown:  $\forall w \exists x (\phi \land dep(v, x) \land var_2(\emptyset, x))$
  - b. <u>free choice</u>:  $\bigvee$  $\forall w \exists x (\phi \land var_{|D|}(v, x))$



## Prove that the generalized variation atom models free choice: